Overview

You will write a program to solve one of the most fundamental NP-Complete problems: determining the satisfiability of a propositional (aka ‘Boolean’) formula presented in conjunctive normal form (CNF). You should use the backtracking-based Davis-Logemann-Loveland algorithm. You should implement unit propagation but you can skip checking for pure literals.

Graduate students will also implement an incomplete model finder — I recommend the WalkSAT/SKC stochastic local search algorithm. For WalkSAT’s \textit{maxFlips} parameter, just hardcode a value that causes about 10 seconds between tries for a largish problem, and for the noise probability \( p \), just use 0.5. For \textit{maxTries}, use something large enough to run for at least a minute.

Input

Your program will read a formula from standard input. Undergraduates won’t take any command-line arguments; graduate students should take \(-w\) to signal use of the WalkSAT algorithm.

We will use a subset of the standard DIMACS CNF format:

c This is a tiny example.
p cnf 4 3
1 -2 3 0
2 3 -4 0 1 4 0

Our format is:

\begin{itemize}
  \item Comments lines start with \texttt{c}.
  \item Before any clauses comes a line starting with \texttt{p cnf}. The rest of the line will be two numbers: the number of variables and the number of clauses.
  \item Clauses are a list of variables, terminated by \texttt{0}. Negative literals have the variable number negated. Note that newlines have no meaning — the \texttt{0} is the clause terminator.
\end{itemize}

Output

We will use the standard DIMACS output format:

\begin{itemize}
  \item After \texttt{s cnf} comes the solution: \texttt{1} for satisfiable, \texttt{0} for unsatisfiable, and \texttt{-1} if the solver couldn’t tell (this would be useful if you implemented a timeout for WalkSAT, for example). After this comes the number of variables and number of clauses (just as in the input \texttt{p} line).
  \item \texttt{v} introduces a variable assignment, where negative means the variable should be set to false.
\end{itemize}

If the formula is not satisfiable, the \texttt{v} lines are omitted.

For DLL, your output should end with the number of branching nodes explored during the search:

\texttt{345 branching nodes explored.}
For WalkSAT, you should list the number of tries and the total number of flips:

14 tries.
2345 total flips.

Execution

Please use make.sh and run.sh scripts as usual.

We supply:

*.cnf a few example benchmark problems.

make-3cnf.py takes two command line arguments: the number of variables and the number of clauses. Outputs a random 3-CNF formula. The critical constrainedness should be around 4.3 clauses per variable.

sat-reference a combined reference solution. Defaults to DLL but runs WalkSAT if given the -w flag.

sat-validator runs your program and validates its output. For example:

```
  sat-validator ./run.sh < tiny-1.cnf
```

or

```
./make-3cnf.py 100 400 | ./sat-validator ./sat-reference -w
```

Thevalidator expects a formula on standard input and will pass it to your program. A one minute timeout is hardcoded into the validator.

Submission

Electronically submit your solution using the instructions on the course web page, including your source code as well as a transcript of your program running with the validator. Include in your write-up answers to the following questions:

1. Describe any implementation choices you made that you felt were important. Clearly explain any aspects of your program that aren’t working. Mention anything else that we should know when evaluating your work.

2. What is the size of the state space for this problem?

3. When there are around 3.3 clauses per variable, how does the largest 3-CNF formula your DLL can solve compare to the largest for your WalkSAT? (Undergrads can run the reference WalkSAT.)

4. Repeat the previous question at 4.3 clauses per variable, and explain the results you obtain.

5. What suggestions do you have for improving this assignment in the future?

Evaluation

Grading is done with an automated script, so be sure you have one algorithm working before starting the next, so that we can give you some credit even if you don’t implement both.

The written questions are one point. Grad students will have 5 points for DLL, 4 points for WalkSAT.

Design Suggestions

It’s better to be slow and correct than fast and buggy.

Just as in a CSP, if you want to be fast, only some of the clauses need to be checked after a variable is assigned or flipped.