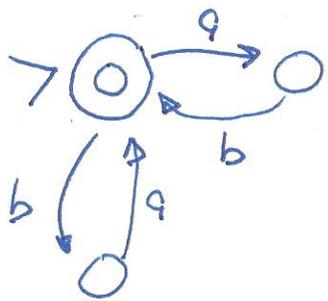


# Finite Automata

CS712

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example



→ the initial state

⊙ a final state

accept start in initial state,  
"consume" the string one symbol  
at a time by traversing an arc  
from the current state labelled  
with the symbol, and when  
string is fully consumed you  
are in a final state.

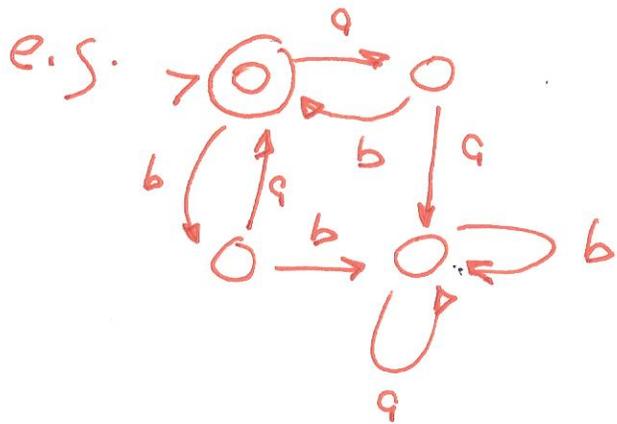
reject end in a non-final state

The language of a finite automaton is the set  
of strings it accepts.

# deterministic finite automaton (DFA)

each state has exactly one src leaving it  
for each alphabet symbol

note usually don't draw "error" srcs  
they all go to a single, non-final  
"sink" state

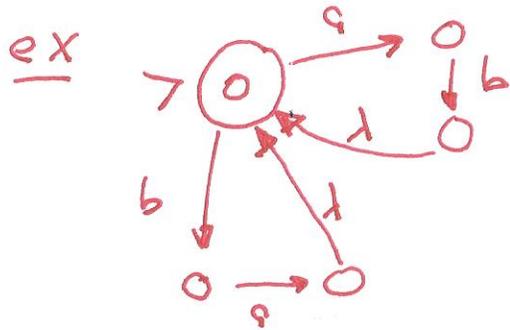


# non deterministic finite automaton (NFA)

a state may have more than one arc  
labelled with the same symbol leaving it

an arc can be labelled with  $\lambda$

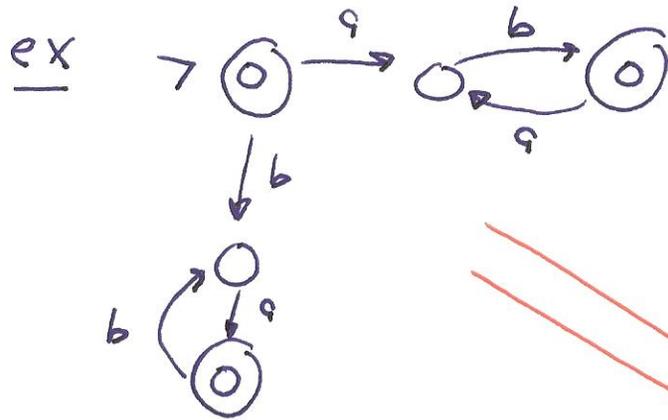
↳ i.e. no input consumed when arc is traversed



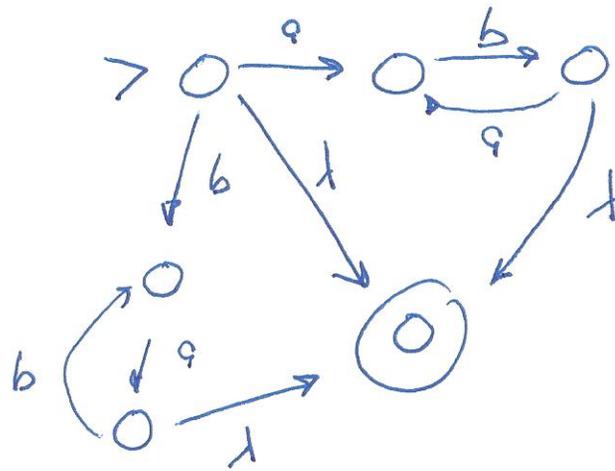
an NFA accepts a string  $w$   
if there is a path from the  
initial state to some final  
state such that  $w$  is the  
concatenation of the symbols  
on the arcs of the path.

otherwise the NFA rejects  
the string.

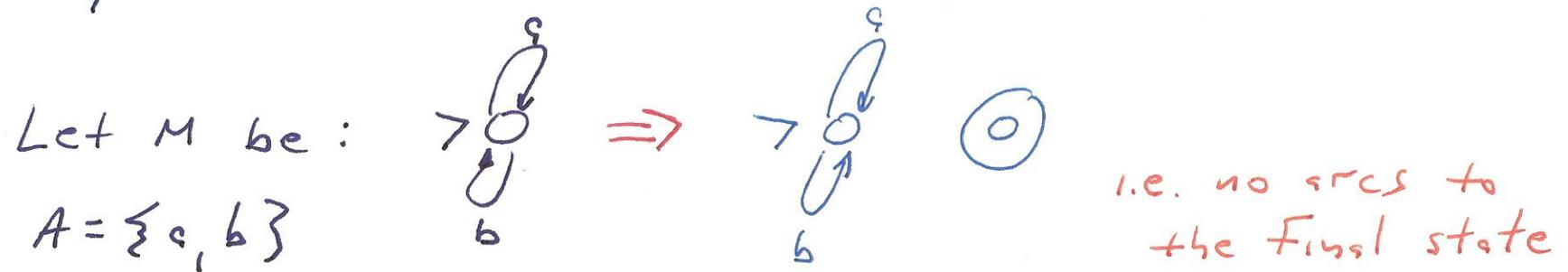
All NFAs with more than one final state can be transformed to equivalent NFAs with a single final state.



$$((ab)^*) \mid ((ba)^*)$$



All NFAs with no final states can be transformed to equivalent NFAs with a single final state.



$$L(M) = \emptyset$$

All NFAs with an initial state that is also a final state can be transformed to equivalent NFAs with initial states that are not final states.

