Thin Locks

Dept. of Computer Science Univ. of New Hampshire If the level of contention is expected to be low, then use lightweight locking mechanism

1.e. use compacting & busy waiting

call yield if can't obtain lock after

some number of tries

CMPX CHG r32, m32

compare eax with m32

if equal, set 2F and store r32 in m32.

else, clear 2F and load m32 into eax.

LMPXCHG must be used with the LOCK prefix
to mike its two memory accesses atomic.

This provides a "thin" lock via the cmpxchg instruction.

```
# Here is the C prototype for the function:
# int thinLock(int *lock, int tryCount);
```

The first parameter is the address of the memory location that is serving as the lock. A zero value in the lock word means the lock is available. A non-zero value means the lock is locked.

The second parameter is a count for how many attempts should be made to obtain the lock before giving up and returning 0 (failure).

The function returns 1 if the lock is obtained and 0 otherwise.

# There is no assembly language thinUnlock because unlock is done by # simply assigning zero to the word that is the lock.

```
.text
                                    # assemble instructions
      .align 4
                                    # put start of function on 4-byte boundary
      .qlobl thinLock
                                    # make function name visible to linker
thinLock:
      pushl
              %ebp
                                    # save old frame pointer
      movl
              %esp, %ebp
                                    # establish new frame pointer
 pusni
              ₹ebx
                                    # save ebx since it is callee saved
      movl
              8(%ebp), %ebx
                                    # get first parameter into ebx
              12(%ebp), %ecx
      movl
                                    # put second parameter into ecx
tryAgain: .
              $0, %eax
      movl
                                    # 0 means lock is available
              $1, %edx
      movl
                                    # put 1 into lock if it is available
      lock
                                    # lock the memory bus for next instruction
      cmpxchg %edx, (%ebx)
                                    # is lock available? (ie (ebx) == eax == 0)
      iе
              gotLock
                                   # if so, done
      subl
              $1, %ecx
                                    # if not, decrement counter
      ie
              qiveUp
                                    # if counter > 0, try again
      gmj
              tryAgain
giveUp:
                                    # eax will be 1 if we branch here
              $0, %eax
     movl
                                   # need to return 0 however
              exit
      jmp
gotLock:
                                   # eax will be 0 if we branch here
     movl
              $1, %eax
                                   # need to return 1 however
exit:
popl >
              %ebx
                                   # restore ebx
     popl
              %ebp
                                   # restore ebp
     ret
                                   # return eax
```