

CS 925

Lecture 5

Traffic Management

Tuesday, February 6, 2024

Little's formula

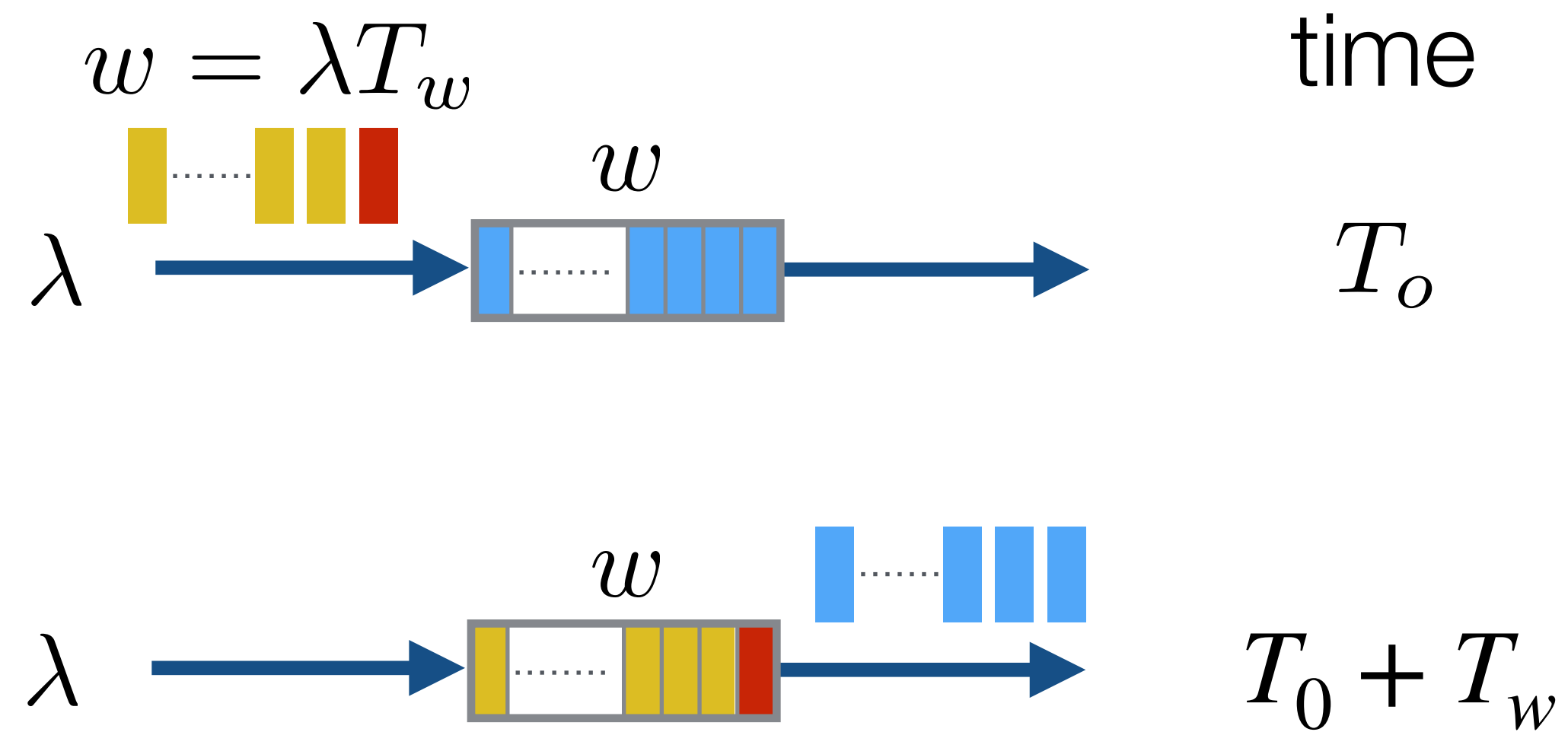
- ▶ **Warning:** lots of handwaving follows

$$w = \lambda T_w$$

- ▶ Assume the system is **stationary** (can be modeled by a stationary process)
 - observing an item: the number of items in the system is, on average, the same when it enters as it is when it leaves
 - it takes T_w from entry to departure, during that period λT_w items arrive (and depart) to maintain the constant number of items in the system
 - assuming FIFO, all items that were in the system prior to the arrival of the observed one must be gone, so the system contains λT_w of the items

Little's formula

- ▶ The question remains: Is the system stationary?
 - under what conditions?
 - does it converge to that state?



Networks of Queues

- ▶ Traffic partitioning and merging, queues in tandem,...
- ▶ **Jackson's Theorem** (1963):
 - Assuming:
 - nodes provide independent service
 - Poisson arrivals from outside
 - fixed partitioning probability
 - no transport delay
 - then, mean delays can be added together
- ▶ ... not really: the theorem does not hold, an error in the proof was found in 2003!

Network Performance

- ▶ **Load vs Latency** diagram
 - impact of load on the delay in delivery
- ▶ **Offered vs Carried Load** diagram
 - impact of load on effective throughput
- ▶ **Loss vs Throughput** diagram
 - impact of packet loss on throughput

Traffic Management

Traffic Management

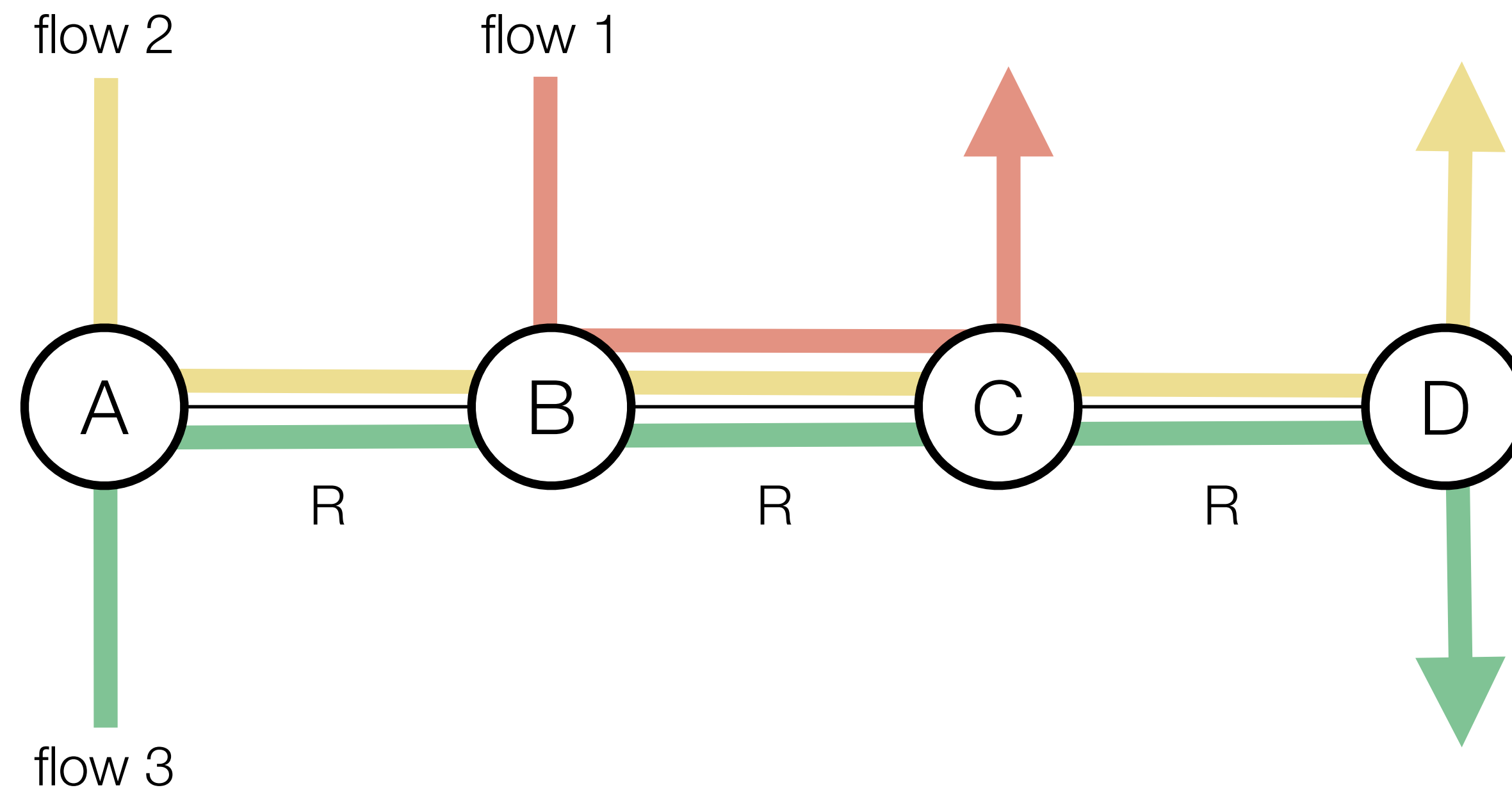
▶ A process to:

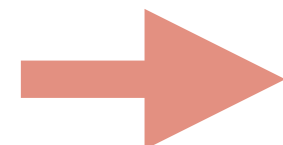

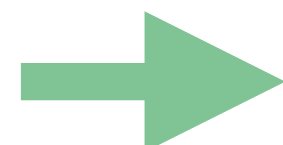
- maximize utility
- ensure fairness
- deliver “quality” service (QoS, QoE, ...)

▶ Where?

- transport (e.g., TCP window)
- network (e.g., obsolete ICMP Source Quench)
- link (e.g., Data Center Bridging)
- application (e.g., HTTP/2, HTTP/3)

Utilization vs fairness



		Max utilization	Max fairness
flow 1		0	$R/3$
flow 2		$R/2$	$R/3$
flow 3		$R/2$	$R/3$

R - link rate

Layers of Traffic Management

- ▶ Within a **device** (router/switch/host)
 - what to do to deliver desired results?
- ▶ Within a **protocol** (protocol layer)
 - how instruct individual devices what they are supposed to do?
- ▶ Within a **network**
 - how to ensure that appropriate level of service is delivered?

Considerations

Obvious but worth reminding ourselves:

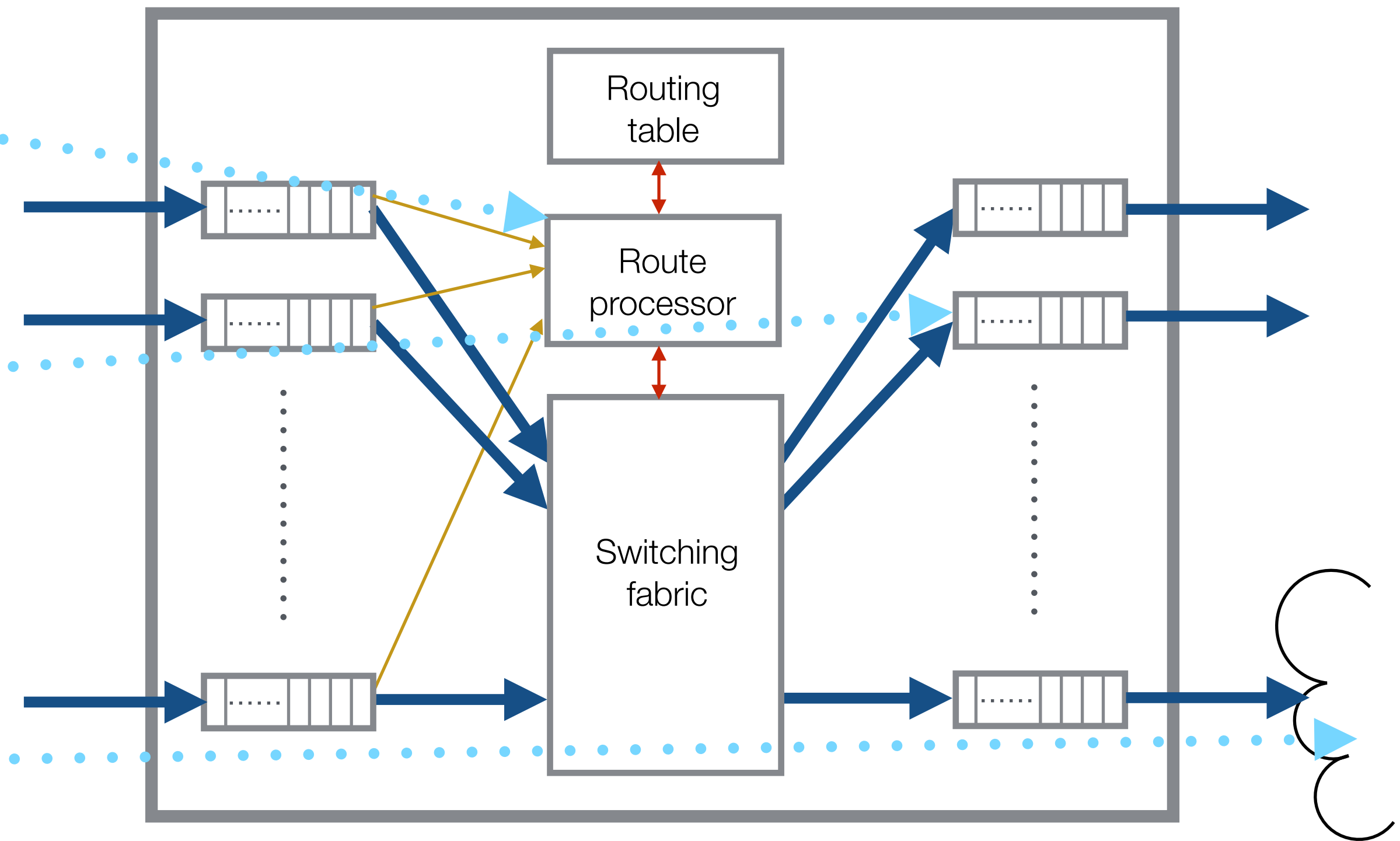
- ▶ In **low load** situations everyone gets the best service possible
 - unless we don't want to create unrealistic performance expectations
- ▶ In **high load** situations, **better service for some** means **degraded service for others**
 - how to determine who “deserves” better service?
 - greater good or more profit?

Questions

- ▶ **How do we know** what to do?
 - methods and techniques that translate user/application demands in traffic management objectives
- ▶ **How do we instruct** the network elements?
 - protocols to facilitate network management information exchange
- ▶ **What do the network elements need to do?**
 - methods and techniques through which the traffic management is implemented in the network

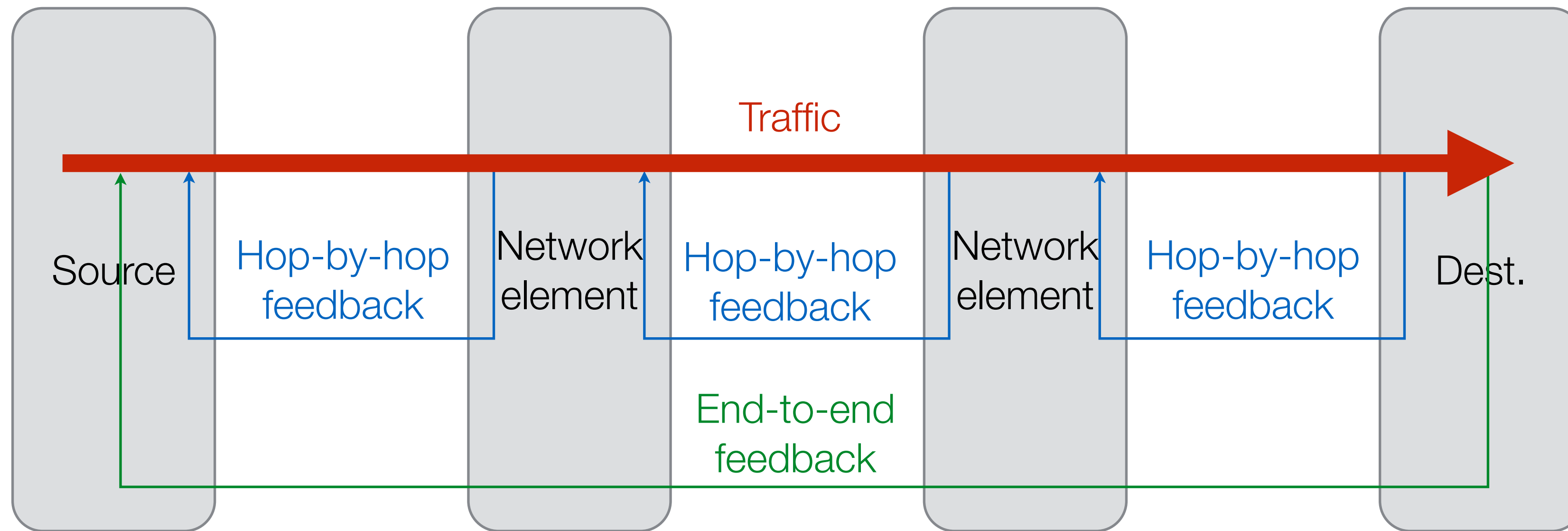
Router/Switch Actions

- ▶ Route selection
- ▶ Queueing policy
- ▶ Subnetwork service request



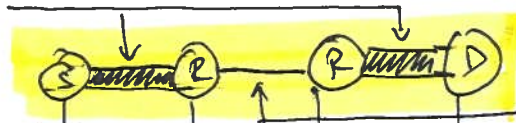
Node to Node Actions

- ▶ Back-pressure / feedback / ...



- ▶ Speed of reaction - feedback loop latency
- ▶ Stability

HIGHER RATE LINKS



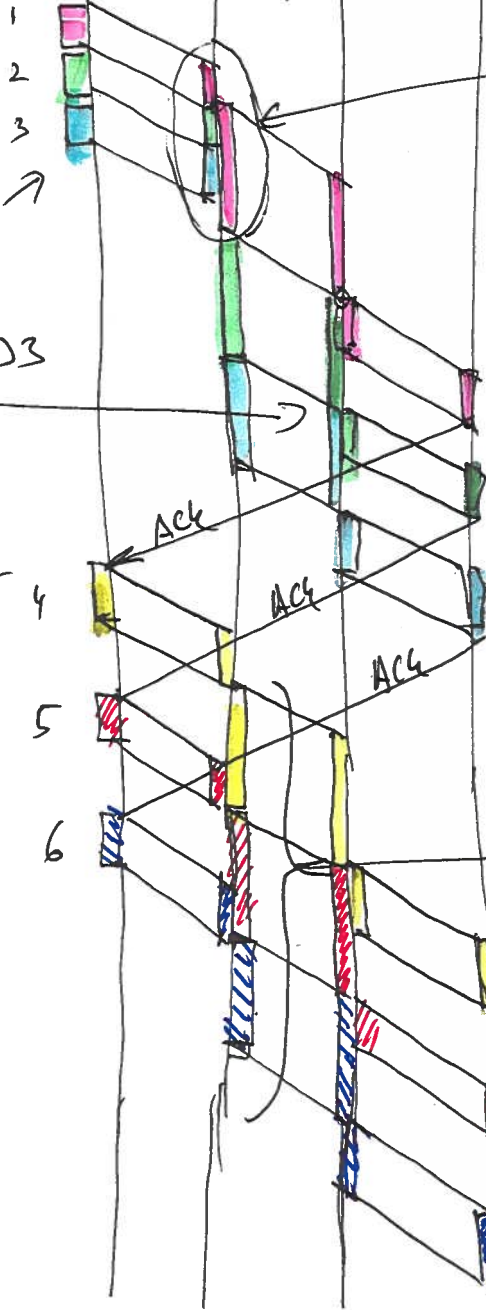
BOTTLE NECK (LOW RATE) LINK

packet rec'd from high rate link but sent over slow link

①

W=3

PACKETS SENT BACK-TO-BACK



②

PACKETS 2 AND 3 QUEUED

③

ACK TRANSN. SPACED BY DATA PACKET ARRIVALS

④

SECOND BATCH SENT WITH GAPS THAT REFLECT THE BOTTLENECK LINK RATE

⑤

PACKETS SPACED, NO NEED FOR QUEUING

