CS 725/825 & T 725Lecture 21 Network Layer

November 20, 2023



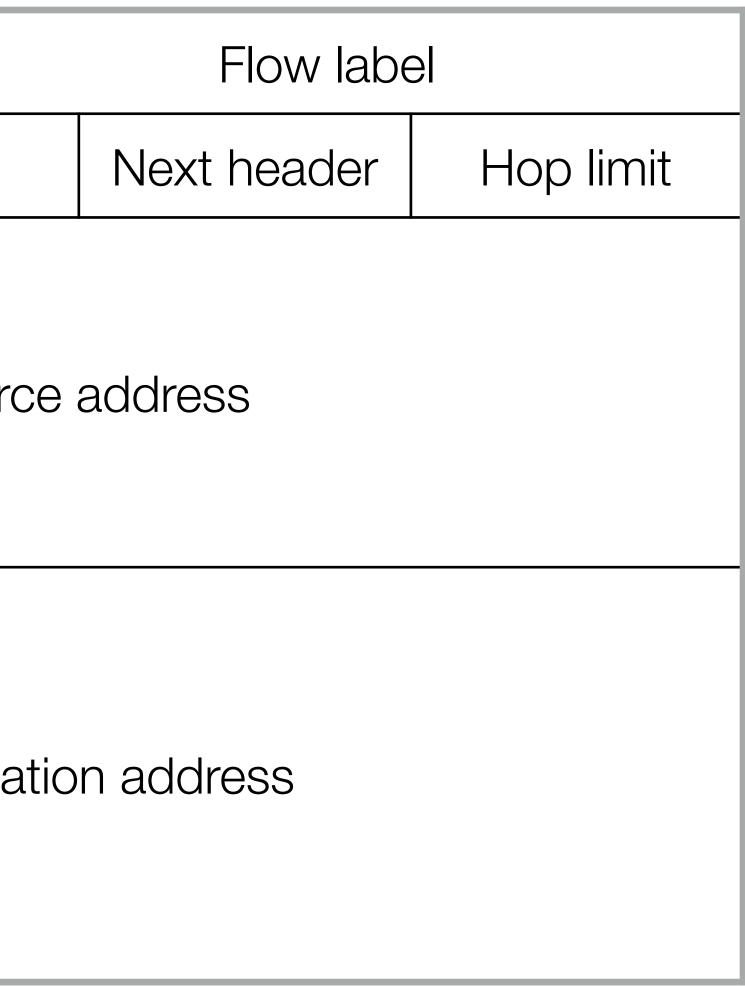
IPv6 - Motivation

- What's wrong with IPv4?
 - not enough addresses
 - to complex to process in routers
 - autoconfiguration
 - security
- Can we avoid switching to IPv6?
 - Network Address Translation (NAT)

IPv6 - Protocol Design

- Keep the good stuff...
 - unreliable datagram service
 - TTL, TOS (for compatibility)
- Eliminate the unnecessary...
 - no fragmentation (only as an option)
 - no header checksums
- Address the issues...
 - longer addresses and more

IPv6 Header



IPv6 Address Representation

- An IPv6 address is represented by 8 groups of 16-bit hexadecimal values separated by colons (:)
- Can be abbreviated:
 - omit leading zeroes in a 16-bit value
 - replace one group of consecutive zeroes by a double colon
- Example:
 - 2606:4100:38c0:9::5 vs 2606:4100:38c0:0009:0000:0000:0000:0005

Special Use IPv6 Addresses

- ::/128 Unspecified address
- ::1/128 Loopback address
- ::FFFF:0:0/96 IPv4-mapped address
- FE80::/10 Link-local unicast
- FF00::/8 Multicast

Routing

- Approaches:
 - Routing)
 - by-hop Forwarding)

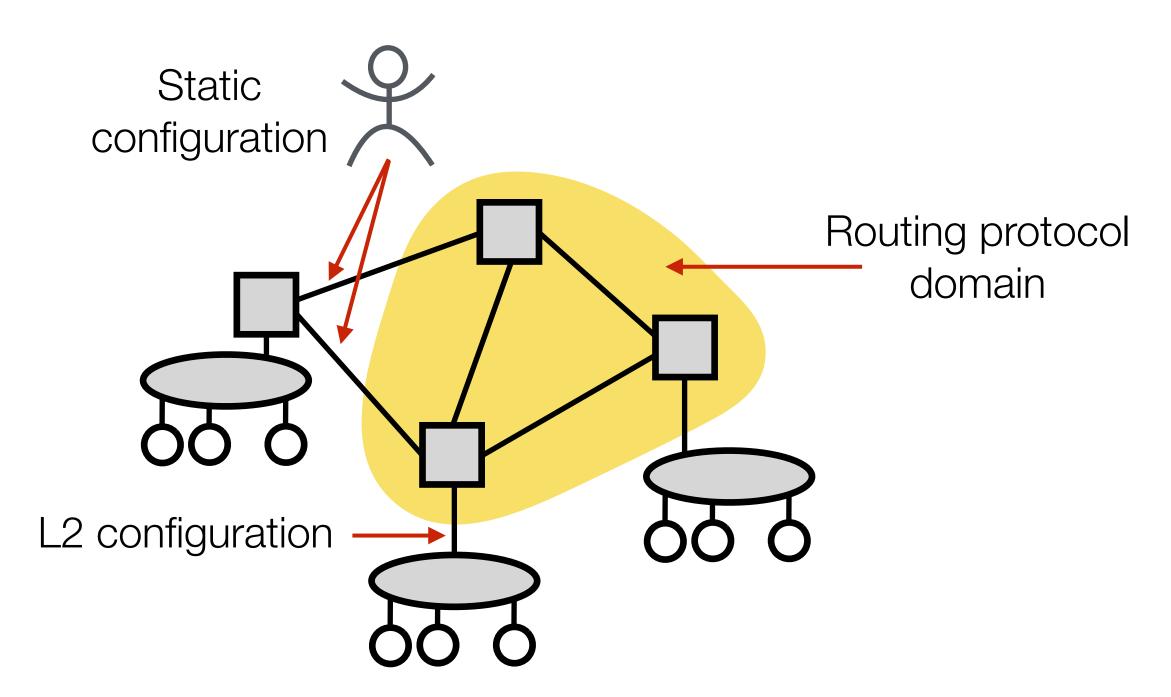
- First find a path from source to destination and then follow it... (Source

- Go to the first corner, ask for direction to the next corner that is on the way to the destination^{*}. Repeat until you reach the destination... (Hop-

* Routing tables give you that information

Routing Table Content

- Automatically populated with entries based in local L2 configuration
- Static entries added by the network administrator Dynamic entries - added by dynamic routing protocols

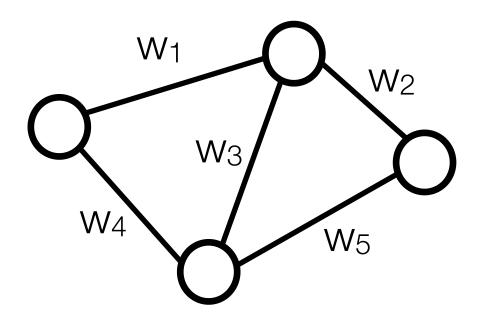


Routing

Finding a good path from source to destination

- topology discovery
- route selection
- Network as a graph...
 - links (point to point and L2 subnets) and routers
 - destinations are typically L2 subnets, not individual nodes
 - links may have "weights"

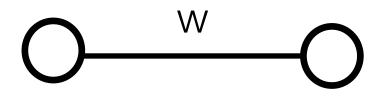
burce to destination



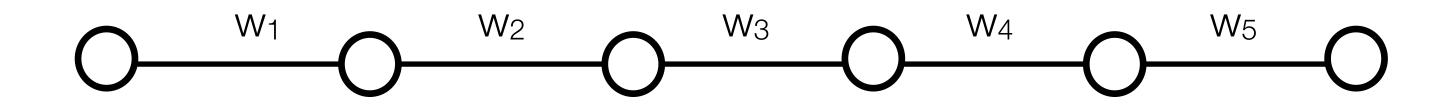


Link weights

What is a good measure of "weight" of a link?



Weight of a path?



Link & Path Measures

- Link measures:
 - Throughput / bit rate
 - Latency
 - Loss probability
 - Availability
 - Current load
 - Security
 - Monetary cost

Path measures:

- Sum

Latency

Monetary cost

- Min/Max

Throughput / bit rate

– Product

Loss probability

Trivial routing methods

- Hot potato routing (not practical)
 - send to randomly chosen outgoing link...
- Flooding (not practical)
 - send a copy to every outgoing link...
- Limited flooding
 - every packet has a sequence number (together with the source address, this makes a copy of a packet uniquely identifiable)
 - send a copy to every other outgoing link
 - keep track of forwarded packets so that copies are sent only once

Routing Protocols - Categories

Link State

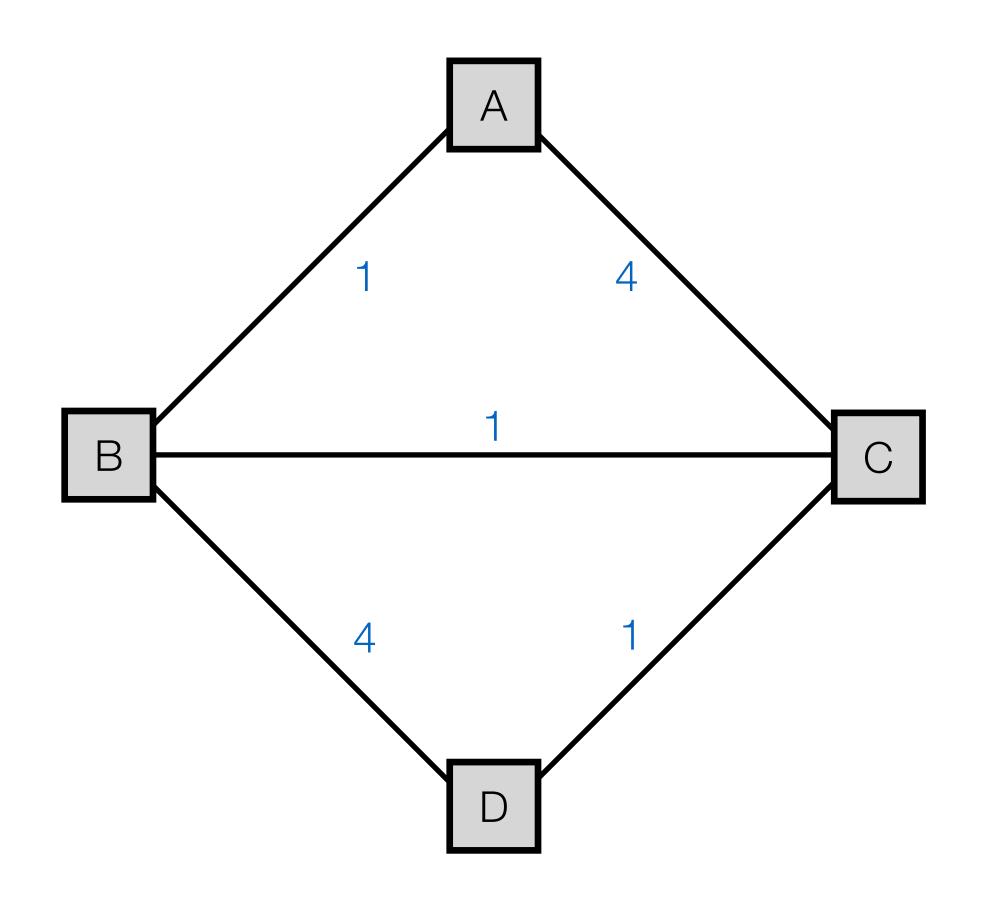
- exact neighbor information flooded to everyone
- topology of the entire networks is discovered in each node
- shortest paths calculated and used to populate the routing tables

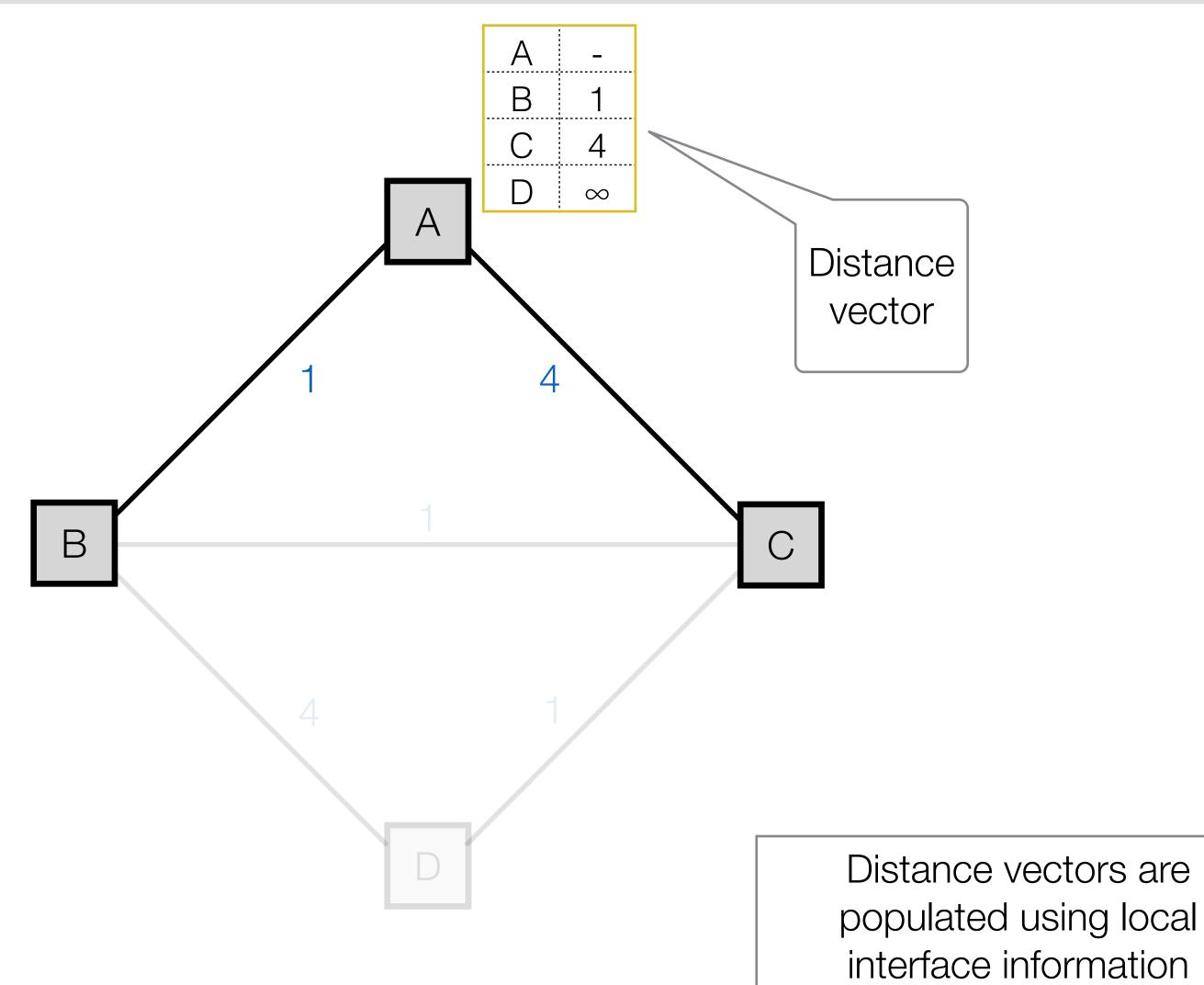
Distance Vector

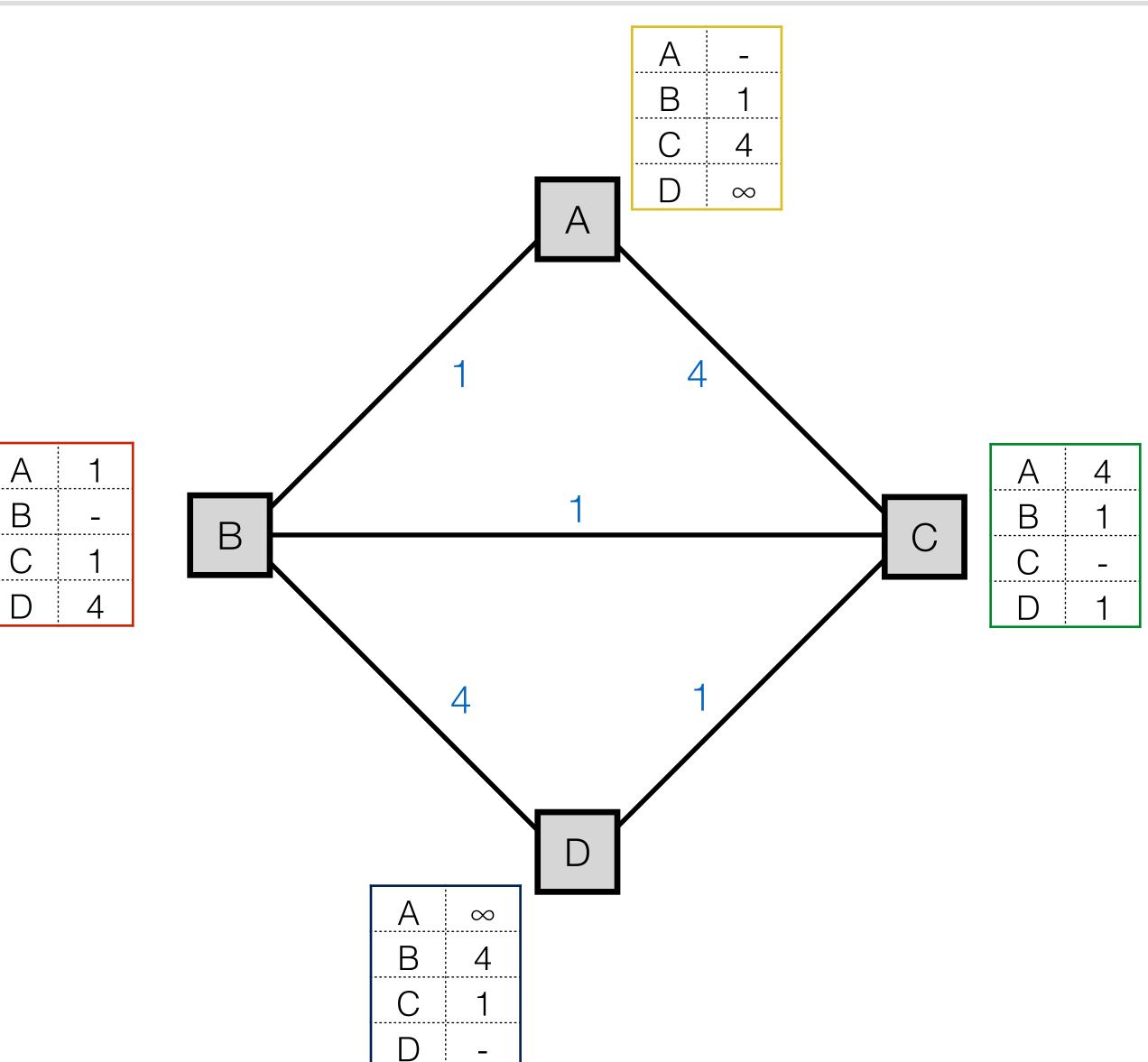
- estimates of distances to all nodes in the network sent to all neighbors
- estimates are improved based on information from neighbors
- the process is repeated and routing tables are populated based on the estimates

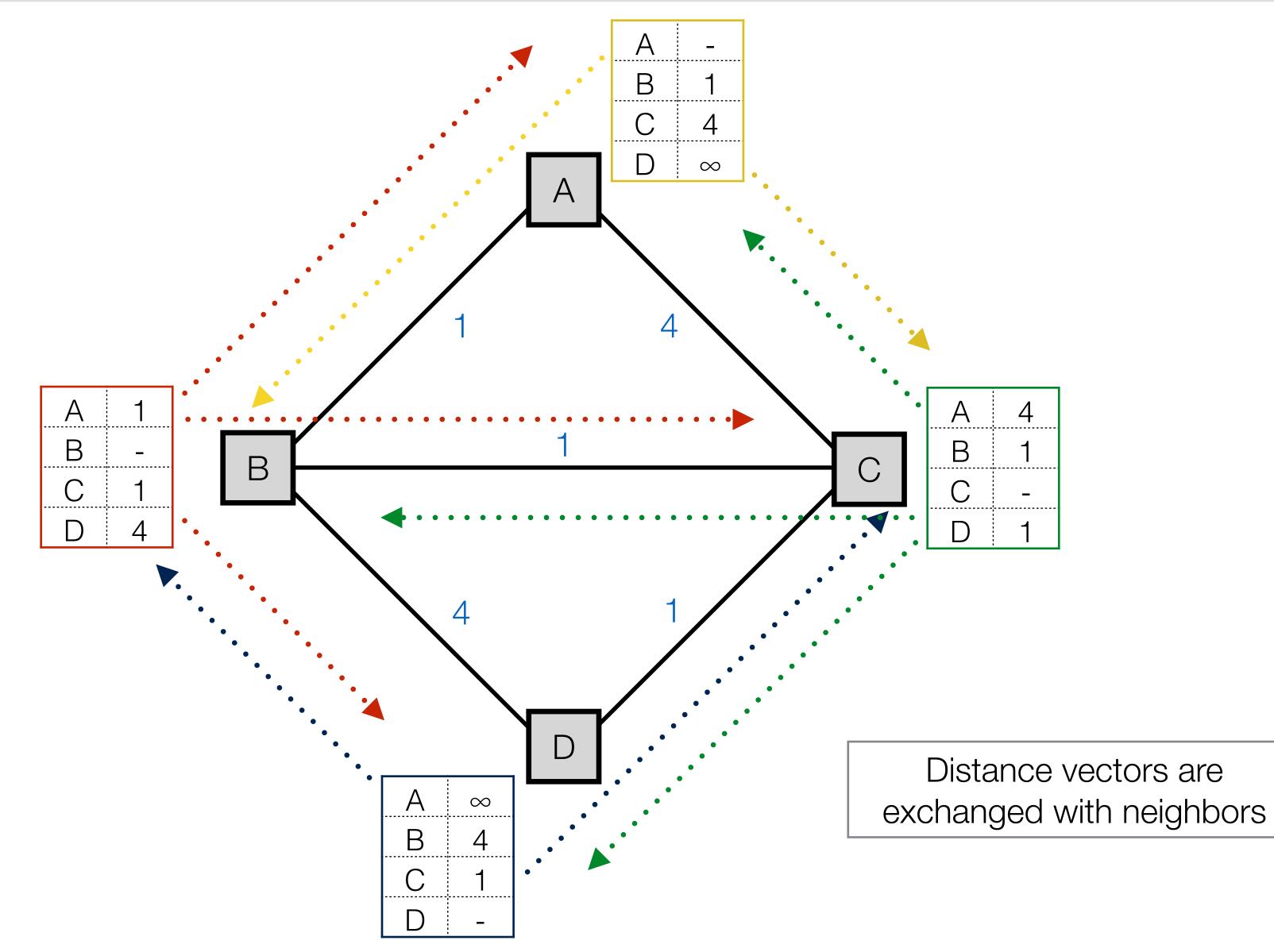
Distance Vector (recap)

- Estimates of distances to all nodes in the network (Distance Vector) is sent to all neighbors
- Estimates are improved based on information from neighbors
- The process is repeated and routing tables are populated based on the estimates

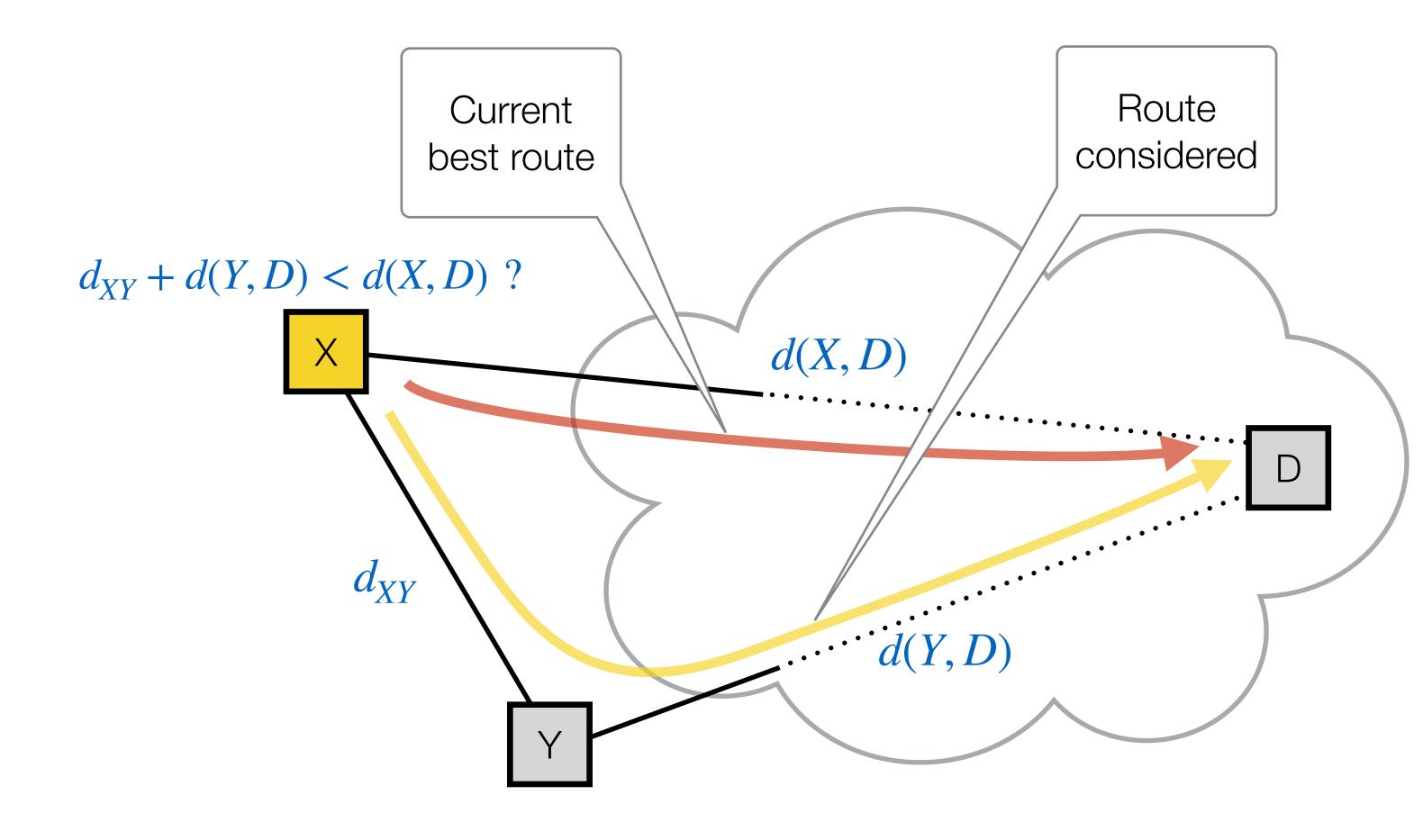




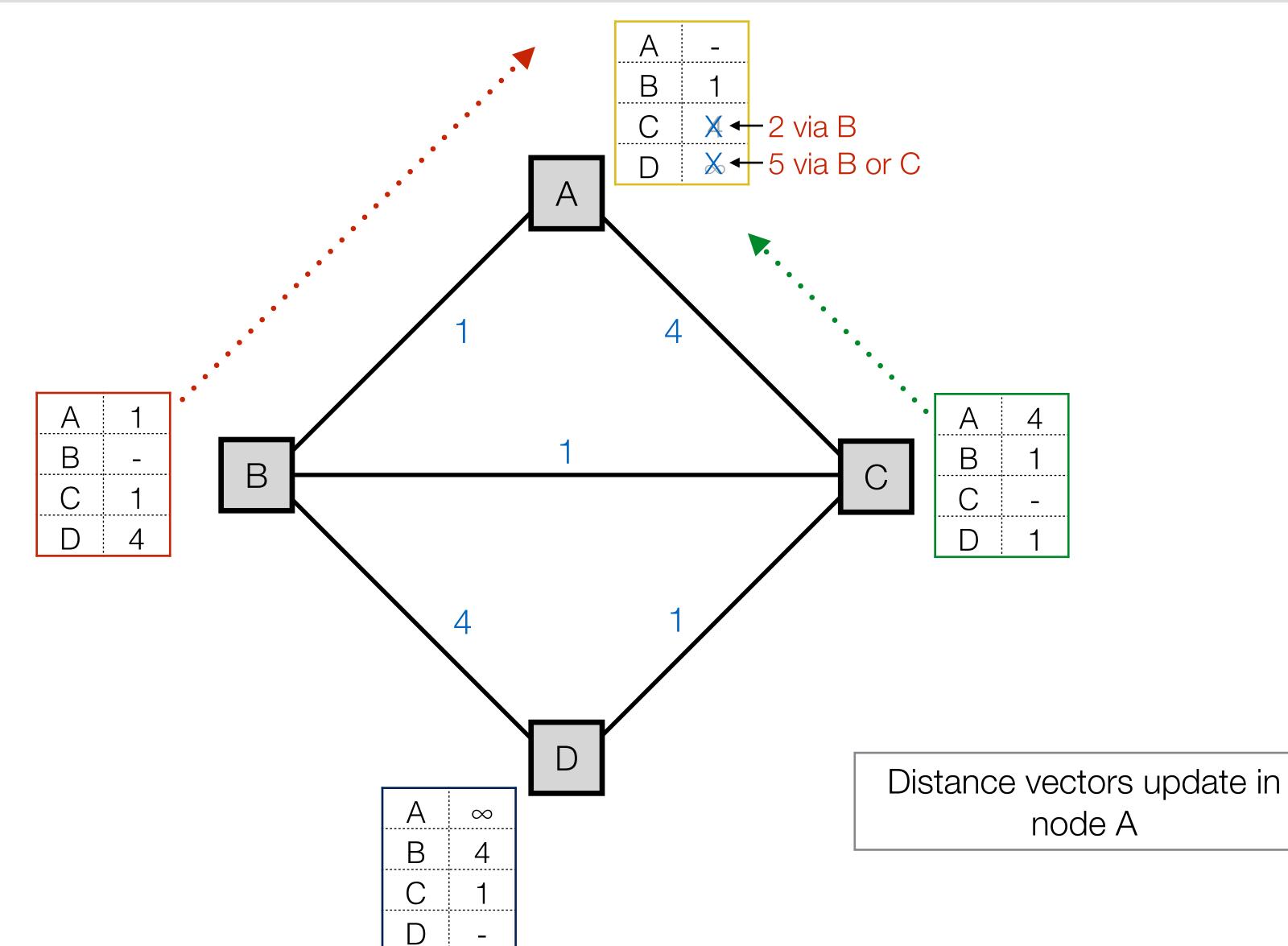


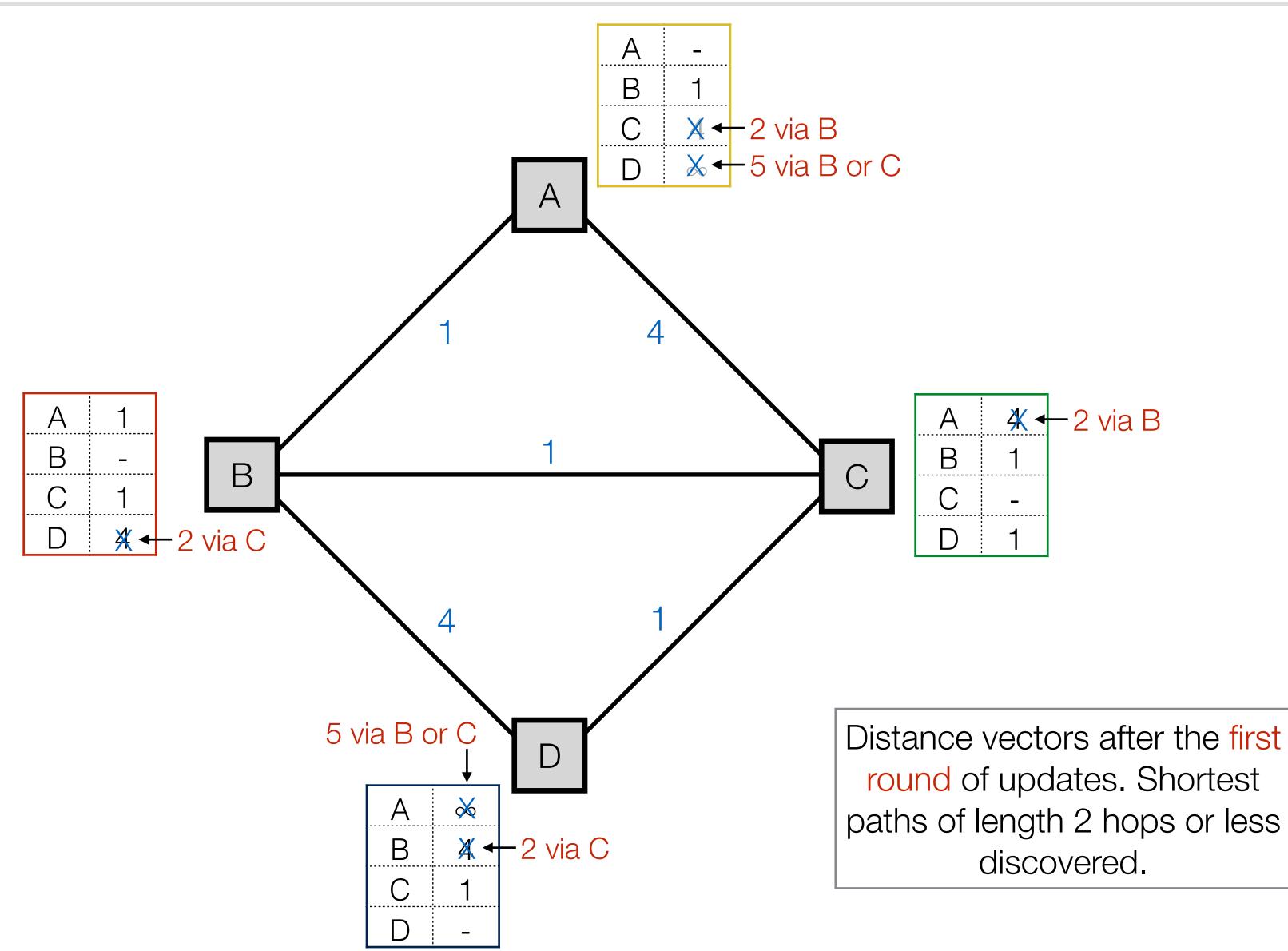


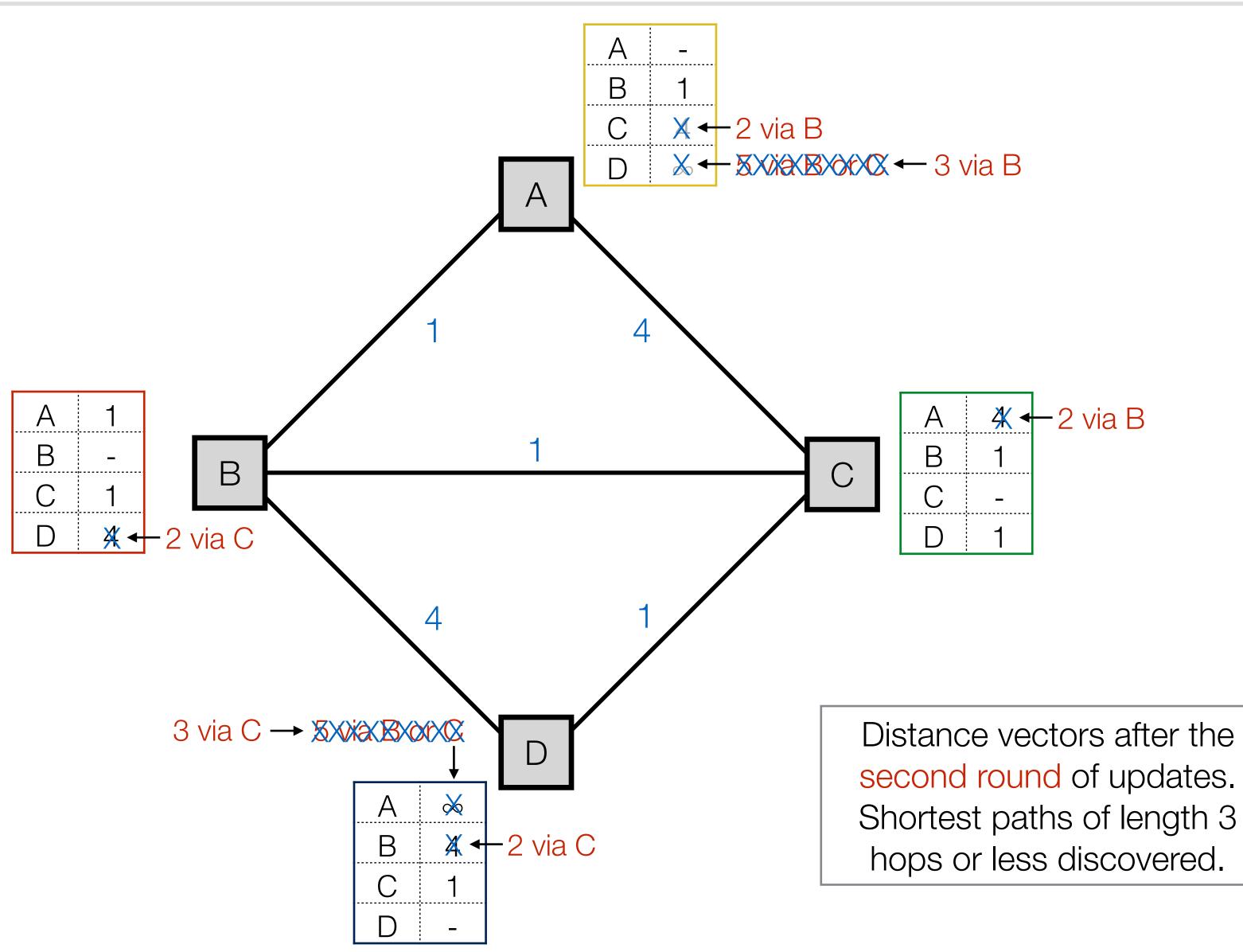
Distance Vector update



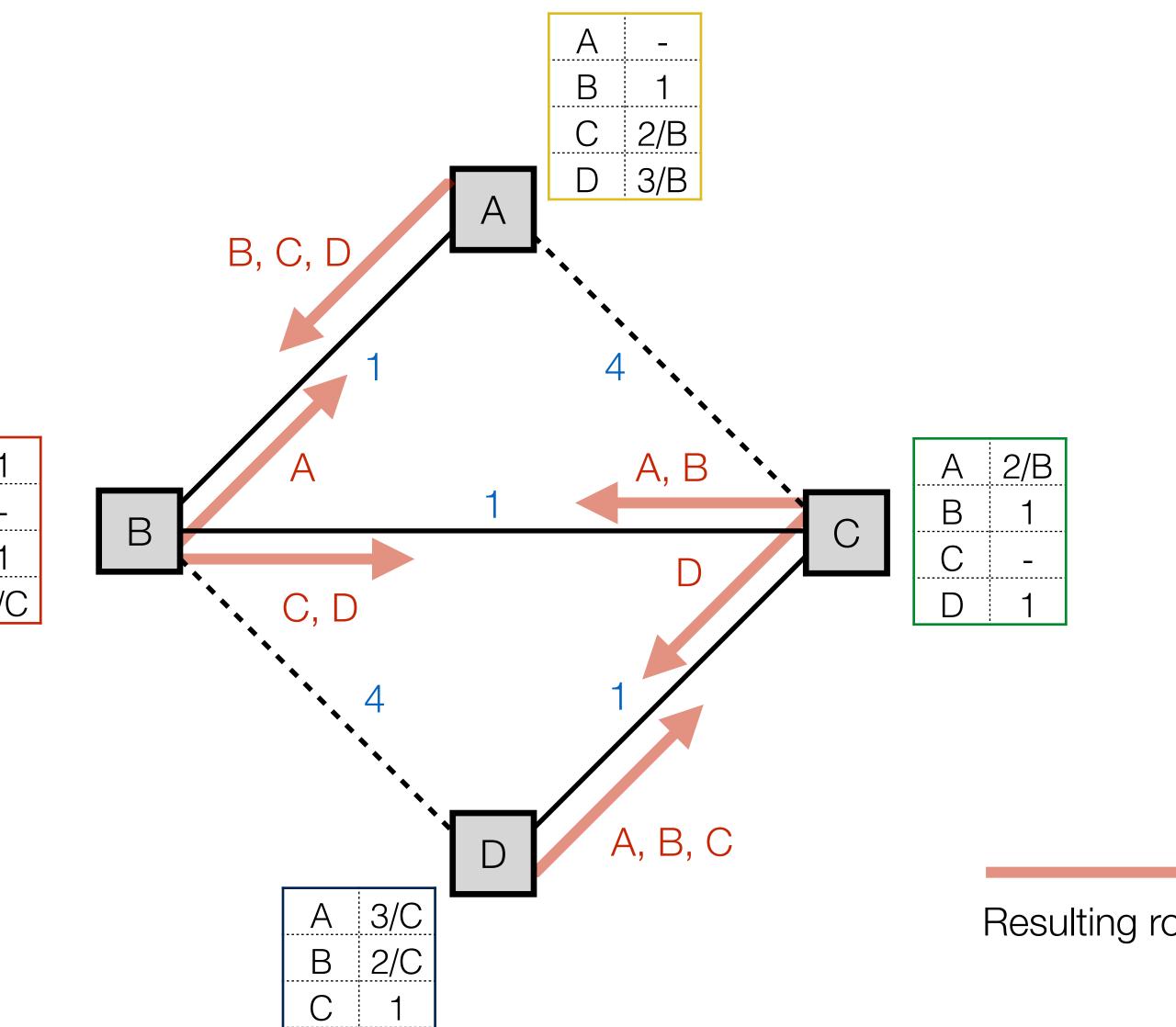
Is routing to D through X's neighbor Y (with distance $d_{XY} + d(Y, D)$) better than the current best route from X (with distance d(X, D))?

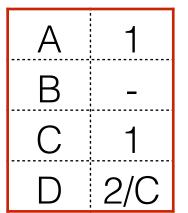






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Resulting routing table content