

# CS 758/858: Algorithms

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<http://www.cs.unh.edu/~ruml/cs758>

[Tries](#)

[Algorithms](#)

[DP](#)

## Tries

- Problem
- Searching
- Searching
- Tries
- Not Tries
- Searching
- Problem
- Break

Algorithms

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DP

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# Tries

# Problem Statement

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Given a linked list of items, print all possible subsets.

running time?

Tries

Problem

Searching

Searching

Tries

Not Tries

Searching

Problem

Break

Algorithms

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DP

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# Searching

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## Tries

■ Problem

■ Searching

■ Searching

■ Tries

■ Not Tries

■ Searching

■ Problem

■ Break

## Algorithms

DP

Structure	Find	Insert	Delete
List (unsorted)			
List (sorted)			
Array (unsorted)			
Array (sorted)			
Heap			
Hash table			
Binary tree (unbalanced)			
Binary tree (balanced)			

# Searching

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## Tries

- Problem
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## Algorithms

## DP

What about long keys?

Can we detect miss without examining entire key?

## Tries

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## Algorithms

## DP

trie: test digits of key, branching on values

- some nodes do not hold values
- fixed order
- depth = length
- canonical representation

retrieval

CLRS: 'trie' = 'radix tree'

Wikipedia: 'trie'  $\neq$  'radix tree'

Sedgewick: 'trie'  $\neq$  'digital search tree'

duplicate keys?

what's their weakness?

# Not Tries

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## Tries

- Problem
- Searching
- Searching
- Tries
- **Not Tries**
- Searching
- Problem
- Break

## Algorithms

## DP

Wikipedia 'radix tree' = 'radix trie' = 'patricia trie': compressed trie, every internal node has at least two leaves beneath

Sedgewick: 'digital search tree': value at every node, just like binary trees except test bits instead of full compare

# Searching

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## Tries

- Problem
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- Break

## Algorithms

## DP

Structure	Find	Insert	Delete
List (unsorted)			
List (sorted)			
Array (unsorted)			
Array (sorted)			
Heap			
Hash table			
Binary tree (unbalanced)			
Binary tree (balanced)			
Trie			



# Problem Statement

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Given a list of records, which may contain duplicates, return a list containing each record at most once.

## Tries

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- Problem
- Searching
- Searching
- Tries
- Not Tries
- Searching
- **Problem**
- Break

## Algorithms

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## DP

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# Break

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## Tries

---

- Problem
- Searching
- Searching
- Tries
- Not Tries
- Searching
- Problem

## ■ Break

## Algorithms

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## DP

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■ asst 4

■ asst 5

Tries

Algorithms

■ Algorithms

■ Types of Algs

DP

# Algorithms

Tries

Algorithms

■ Algorithms

■ Types of Algs

DP

Beyond craftsmanship lies invention, and it is here that lean, spare, fast programs are born. Almost always these are the result of strategic breakthrough rather than tactical cleverness. Sometimes the strategic breakthrough will be a new algorithm, such as the Cooley-Tukey Fast Fourier transform or the substitution of an  $n \log n$  sort for an  $n^2$  set of comparisons.

Much more often, strategic breakthrough will come from redoing the representation of the data or tables. This is where the heart of a program lies. Show me your flowchart and conceal your tables, and I shall continue to be mystified. Show me your tables, and I won't usually need your flowchart; it'll be obvious.

— Fred Brooks, 1974 (lead on IBM System/360, Turing Award 1999)

# Algorithms: Modern Version

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Tries

Algorithms

■ Algorithms

■ Types of Algs

DP

Smart data structures and dumb code works a lot better than the other way around.

— Guy Steele, 2002 (ACM Fellow, inventor of Scheme, editor of *The Hacker's Dictionary*)

# Types of Algorithms

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Tries

Algorithms

■ Algorithms

■ Types of Algs

DP

- divide and conquer
- dynamic programming
- greedy
- backtracking
- (reduction)

Tries

Algorithms

DP

- Fibonacci
- Memoization
- EOLQs

# Dynamic Programming

# Fibonacci Numbers

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Tries

Algorithms

DP

■ Fibonacci

■ Memoization

■ EOLQs

$$F_n = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ F_{n-1} + F_{n-2} & \text{for } n \geq 2 \end{cases}$$

What is the complexity of the naive algorithm?  
How to make this efficient?



# Memoization

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Tries

Algorithms

DP

■ Fibonacci

■ Memoization

■ EOLQs

- recursive decomposition
- polynomial number of subproblems
- cache results in look-up table

one form of *dynamic programming*

Tries

Algorithms

DP

■ Fibonacci

■ Memoization

■ EOLQs

- What's still confusing?
- What question didn't you get to ask today?
- What would you like to hear more about?

Please write down your most pressing question about algorithms and put it in the box on your way out.

*Thanks!*